# Randomized Numerical Linear Algebra Lecture 1 

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## General info

- Course Website: https://fabianfaulstich.github.io/MATH6950_2024/
- Homework assignments:

1. Submission through Gradescope.
2. Gradescope will close on the due date at 11:59 p.m. No late submissions!
3. Everyone get one joker
4. zero-tolerance policy regarding cheating

- Lectures:

1. hybrid slides and blackboard lecturing
2. the slide part will be made available online
3. Code presented and used in class will be made available online

- Programming assignments:

1. P.A.s will carry substantial points in each homework
2. To gain full credit you have to show exploration and clear reasoning

## Planned course outlook

1. Review of numerical linear algebra and probability theory
2. Low-rank approximations and randomness

## What to expect

- The course is centered primarily on computational aspects
- Aimed at equipping you with a robust set of computational skills
- Proofs will be included (mostly) at a high level
- The course aims to impart an understanding of the key ideas behind proofs rather than delving into exhaustive, fully worked-out proofs


## numerical linear algebra

- solving dense and sparse linear systems



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- solving dense and sparse linear systems
- orthogonalization, least square \& (Tikhonov) regularization
- determination of eigenvalues \& eigenvectors, invariant subspaces
- singular value decomposition (SVD)


$$
M=U \cdot \Sigma \cdot V^{*}
$$

## Numerical linear algebra

What is the problem?
$\Rightarrow$ Large scaling problems!

Example:
Given an $m \times n$ matrix $\mathbf{A}$ where both $m$ and $n$ are large, the singular value decomposition will require memory and time which is superlinear in $m$ and $n$

$$
\begin{equation*}
\mathcal{O}\left(4 m n^{2}-\frac{4}{3} n^{3}\right) \tag{GK}
\end{equation*}
$$

Randomized algorithms bring this down to $m$ and $k$, where $k$ is the rank

$$
\begin{equation*}
\mathcal{O}\left(2 k n^{2}+2 n^{3}\right) \tag{LHC}
\end{equation*}
$$

## Randomized algorithms

1. Monte Carlo (M.C.) algorithms (~1940)

- Last resort methods
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- Hesitation: Two runs should produce the same number


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2. Randomized algorithms in numerical linear algebra ( $\sim 1980$ )

- Power method, random initialization (Dixon 1983)
- M. C. methods for trace estimation (Girard 1989 \& Hutchinson 1990)
- Randomized transformations can avoid pivoting steps in Gaussian elimination (Parker 1995)


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3. Practical randomized algorithms for low-rank matrix approximation and least-squares problems (mid-2000s)

- First computational evidence that randomized algorithms outperform classical NLA algorithms for particular classes of problems


## Review: Classical Numerical Linear Algebra

- L. N. Trefethen and D. Bau III, Numerical linear algebra, Vol. 50, SIAM
- G. Stewart, Matrix Algorithms Volume 1: Basic Decompositions, SIAM
- G. W. Stewart, Matrix algorithms volume 2: eigensystems, Vol. 2, SIAM
- G. H. Golub and C. F. Van Loan, Matrix computations
- R. A. Horn and C. R. Johnson, Matrix analysis
- R. Bhatia, Matrix analysis, Vol. 169 of Graduate Texts in Mathematics


## Basics - Notation I

- Algebraic field: $\mathbb{C}, \mathbb{R}, \mathbb{F}$
- Scalars: $a, b, \ldots$ or $\alpha, \beta, \ldots$
- Vectors are elements of $\mathbb{F}^{n}, n \in \mathbb{N}$ : $\mathbf{a}, \mathbf{b}, \ldots$ or $\boldsymbol{\alpha}, \boldsymbol{\beta}, \ldots$
- Special vectors $\mathbf{0}, \mathbf{1}, \boldsymbol{\delta}_{i} \in \mathbb{F}^{n}$ :

$$
\mathbf{0}=\left(\begin{array}{c}
0 \\
0 \\
\vdots \\
0
\end{array}\right) \quad \mathbf{1}=\left(\begin{array}{c}
1 \\
1 \\
\vdots \\
1
\end{array}\right) \quad \boldsymbol{\delta}_{i}=\left(\begin{array}{c}
0 \\
\vdots \\
0 \\
1 \\
0 \\
\vdots \\
0
\end{array}\right)
$$

- Vector element:

$$
(\mathbf{a})_{i}=\mathbf{a}(i) \quad i \text { th coordinate of } \mathbf{a}
$$

- Colon notation: $(\mathbf{a})_{1: i}=\mathbf{a}(1: i)$


## Basics - Notation II

- A matrix is an element of $\mathbb{F}^{m \times n}, m, n \in \mathbb{N}$ :
$\mathbf{A}, \mathbf{B}, \ldots$ or $\boldsymbol{\Lambda}, \boldsymbol{\Delta}, \ldots$
- Matrix element:

$$
(\mathbf{A})_{i j}=\mathbf{A}(i j) \quad(i, j) \text { th coordinate of } \mathbf{A}
$$

- Special matrices $\mathbf{0}, \mathbf{I} \in \mathbb{F}^{m \times n}$ :

$$
\mathbf{0}=\left(\begin{array}{ccc}
0 & \cdots & 0 \\
\vdots & \ddots & \vdots \\
0 & \cdots & 0
\end{array}\right) \quad \mathbf{I}=\left(\begin{array}{cccc}
1 & 0 & \cdots & 0 \\
0 & \ddots & \ddots & \vdots \\
\vdots & \ddots & 1 & 0 \\
0 & \cdots & 0 & 1
\end{array}\right)
$$

- Colon notation: $(\mathbf{A})_{i:} \equiv i$ th row and $(\mathbf{A})_{: j} j$ th column of $\mathbf{A}$


## Basics - Notation III

- $*$ is the conjugate transpose
- $\mathbb{H}_{n}=\mathbb{H}_{n}(\mathbb{F})=\left\{\mathbf{A} \in \mathbb{F}^{n \times n}: \mathbf{A}=\mathbf{A}^{*}\right\}$
- $\dagger$ is the Moor-Penrose (pseudo)inverse:


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- $\dagger$ is the Moor-Penrose (pseudo)inverse:
$\mathbf{A}^{\dagger}$ is the MP inverse of $\mathbf{A}$ iff

$$
\begin{array}{rlrl}
\text { (i) } & \mathbf{A} \mathbf{A}^{\dagger} \mathbf{A} & =\mathbf{A} & (i i i) \\
(i i) & \left.\mathbf{A}^{\dagger}\right)^{*} & =\mathbf{A} \mathbf{A}^{\dagger} \\
\mathbf{A}^{\dagger} \mathbf{A} \mathbf{A}^{\dagger} & =\mathbf{A}^{\dagger} & & \text { (iv) } \\
\left(\mathbf{A}^{\dagger} \mathbf{A}\right)^{*} & =\mathbf{A}^{\dagger} \mathbf{A}
\end{array}
$$

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\left(\mathbf{A}^{\dagger} \mathbf{A}\right)^{*} & =\mathbf{A}^{\dagger} \mathbf{A}
\end{array}
$$

If $\mathbf{A}$ has full column rank, then $\mathbf{A}^{\dagger}=\left(\mathbf{A}^{*} \mathbf{A}\right)^{-1} \mathbf{A}^{*}$
If $\mathbf{A}$ attains an inverse then

$$
\mathbf{A}^{\dagger}=\mathbf{A}^{-1}
$$

## Eigenvalues and singular values

- $\mathrm{PSD}=\left\{\mathbf{A} \in \mathbb{H}_{n}: \mathbf{x}^{\top} \mathbf{A} \mathbf{x} \geq 0\right.$ for $\left.\mathbf{x} \neq 0\right\}$
- $\mathrm{PD}=\left\{\mathbf{A} \in \mathbb{H}_{n}: \mathbf{x}^{\top} \mathbf{A} \mathbf{x}>0\right.$ for $\left.\mathbf{x} \neq 0\right\}$
- $\preccurlyeq$ denotes the semidefinte order on $\mathbb{H}_{n}$, i.e.

$$
\mathbf{A} \preccurlyeq \mathbf{B} \Leftrightarrow 0 \preccurlyeq \mathbf{B}-\mathbf{A}
$$

- $\prec$ denotes the definte order on $\mathbb{H}_{n}$, i.e.

$$
\mathbf{A} \prec \mathbf{B} \Leftrightarrow 0 \prec \mathbf{B}-\mathbf{A}
$$

- Eigenvalues of $\mathbf{A} \in \mathbb{H}_{n}: \lambda_{1} \geq \lambda_{2} \geq \ldots$
- Singular values of $\mathbf{A} \in \mathbb{F}^{m \times n}: \sigma_{1} \geq \sigma_{2} \geq \ldots$
- Let $f: \mathbb{R} \rightarrow \mathbb{R}$. We extend $f$ to spectral function $f: \mathbb{H}_{n} \rightarrow \mathbb{H}_{n}$

$$
f(\mathbf{A}):=\sum_{i=1}^{n} f\left(\lambda_{i}\right) \mathbf{u}_{i} \mathbf{u}_{i}^{*} \quad \text { where } \quad \mathbf{A}=\sum_{i=1}^{n} \lambda_{i} \mathbf{u}_{i} \mathbf{u}_{i}^{*}
$$

## Inner products and geometry I

- Equip $\mathbb{F}^{n}$ with standard scalar product and associated $\ell^{2}$-norm. Let $\mathbf{a}, \mathbf{b} \in \mathbb{F}^{n}$ then

$$
\langle\mathbf{a}, \mathbf{b}\rangle:=\mathbf{a} \cdot \mathbf{b}=\mathbf{a}^{*} \mathbf{b}=\sum_{i=1}^{n}(\mathbf{a})_{i}^{*}(\mathbf{b})_{i}
$$

and

$$
\|\mathbf{a}\|^{2}:=\langle\mathbf{a}, \mathbf{a}\rangle
$$

- Unit sphere in $\mathbb{F}^{n}: \mathbb{S}^{n-1}=\mathbb{S}^{n-1}(\mathbb{F})$


## Inner products and geometry II

- The trace of $\mathbf{A} \in \mathbb{F}^{n \times n}$ :

$$
\operatorname{Tr}(\mathbf{A})=\operatorname{trace}(\mathbf{A})=\sum_{i=1}^{n}(\mathbf{A})_{i i}
$$

Nonlinear functions bind before the trace.

- Equip $\mathbb{F}^{m \times n}$ with the standard trace inner product and Frobenius norm:
Let $\mathbf{A}, \mathbf{B} \in \mathbb{F}^{m \times n}$ then

$$
\langle\mathbf{A}, \mathbf{B}\rangle:=\operatorname{Tr}\left(\mathbf{A}^{*} \mathbf{B}\right)
$$

and

$$
\|\mathbf{A}\|_{F}^{2}=\langle\mathbf{A}, \mathbf{A}\rangle
$$

- $\mathbf{U} \in \mathbb{F}^{m \times n}$ is orthonormal iff $\mathbf{U}^{*} \mathbf{U}=\mathbf{I}_{n}$.
$\mathbf{U}$ is unitary $(\mathbb{F}=\mathbb{C})$ or orthogonal $(\mathbb{F}=\mathbb{R})$ if $m=n$.


## Matrix norms I

- Let $\|\cdot\|_{\alpha}$ be a norm on $\mathbb{F}^{n}$ and $\|\cdot\|_{\beta}$ be a norm on $\mathbb{F}^{m}$. Then

$$
\|\cdot\|_{\alpha, \beta}: \mathbb{F}^{m \times n} \rightarrow \mathbb{R} ; \mathbf{A} \mapsto \sup _{\substack{\mathbf{x} \in \mathbb{F}^{n} \\\|\mathbf{x}\|_{\alpha} \neq 0}} \frac{\|\mathbf{A} \mathbf{x}\|_{\beta}}{\|\mathbf{x}\|_{\alpha}}
$$

Induces a norm on $\mathbb{F}^{m \times n}$.

- Alternatively, we may define any function

$$
\|\cdot\|: \mathbb{F}^{m \times n} \rightarrow \mathbb{R}
$$

that fulfills:

1. $0 \leq\|\mathbf{A}\|, \forall \mathbf{A} \in \mathbb{F}^{m \times n}$ and $\|\mathbf{A}\|=0 \Leftrightarrow \mathbf{A}=\mathbf{0}$
2. $\|a \mathbf{A}\|=|a|\|\mathbf{A}\|, \forall \mathbf{A} \in \mathbb{F}^{m \times n}$, and $\forall a \in \mathbb{F}$
3. $\|\mathbf{A}+\mathbf{B}\| \leq\|\mathbf{A}\|+\|\mathbf{B}\|, \forall \mathbf{A}, \mathbf{B} \in \mathbb{F}^{m \times n}$

## Matrix norms II

Several matrix norms will be used. Let $\mathbf{A} \in \mathbb{F}^{m \times n}$

- The unadorned norm $\|\cdot\|$ is the spectral norm

$$
\|\mathbf{A}\|=\sigma_{1}=\|\mathbf{A}\|_{\ell^{2}}
$$

- $\|\cdot\|_{*}$ is the nuclear/trace norm

$$
\|\mathbf{A}\|_{*}=\sum_{k=1}^{\min (m, n)} \sigma_{k}
$$

- $\|\cdot\|_{F}$ is the Frobenius norm

$$
\|\mathbf{A}\|_{F}^{2}=\sum_{i=1}^{m} \sum_{j=1}^{n}\left|(A)_{i j}\right|^{2}=\sum_{k=1}^{\min (m, n)} \sigma_{k}^{2}=\operatorname{Tr}\left(\mathbf{A}^{*} \mathbf{A}\right)
$$

## Matrix norms III

- $\|\cdot\|_{p}$ is the Schatten p-norm for $p \in[1, \infty]$

$$
\|\mathbf{A}\|_{p}=\left(\sum_{k=1}^{\min (m, n)} \sigma_{k}^{p}\right)^{\frac{1}{p}}
$$

- $\|\cdot\|_{K, p}$ is the Ky Fan $p$-norm for $p \leq \min (m, n)$

$$
\|\mathbf{A}\|_{K, p}=\sum_{k=1}^{p} \sigma_{k}
$$

Note:

$$
\begin{gathered}
\|\cdot\|_{*}=\|\cdot\|_{K, \min (m, n)}=\|\cdot\|_{1} \\
\|\cdot\|_{F}=\|\cdot\|_{2} \\
\|\cdot\|=\|\cdot\|_{K, 1}=\|\cdot\|_{\infty}
\end{gathered}
$$

## Intrinsic Dimension

Let $\mathbf{A} \in \mathbb{H}_{n}$ be PSD. We define the intrinsic dimension as

$$
\operatorname{intdim}(\mathbf{A}):=\frac{\operatorname{Tr}(\mathbf{A})}{\|\mathbf{A}\|}
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Note that for A non-zero:

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1 \leq \operatorname{intdim}(\mathbf{A}) \leq \operatorname{rank}(\mathbf{A})
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The upper bound is saturated if $\mathbf{A}$ is an orthogonal projector i.e.

1. $\mathbf{A} \in \mathbb{F}^{m \times m}$ and $\mathbf{A}^{2}=\mathbf{A}$
2. $\mathbf{A}$ is projector and $\mathbf{A} \in \mathbb{H}_{n}$

The intrinsic rank can be interpreted as a continuous measure of the rank

## Stable rank

Let $\mathbf{B} \in \mathbb{F}^{m \times n}$. We define the stable rank as

$$
\operatorname{srank}(\mathbf{A}):=\operatorname{intdim}\left(\mathbf{B}^{*} \mathbf{B}\right)=\frac{\|\mathbf{B}\|_{F}^{2}}{\|\mathbf{B}\|^{2}}
$$

## Schur complement

Let

$$
\mathbf{M}=\left(\begin{array}{cc}
\mathbf{A} & \mathbf{B} \\
\mathbf{C} & \mathbf{D}
\end{array}\right) \in \mathbb{F}^{m \times m}
$$

with $\mathbf{A} \in \mathbb{F}^{n \times n}$.
If $\mathbf{D}$ is invertible the Schur complement of $\mathbf{D}$ in $\mathbf{M}$ is

$$
\mathbf{M} / \mathbf{D}:=\mathbf{A}-\mathbf{B D}^{-1} \mathbf{C}
$$

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$$
\mathbf{M} / \mathbf{A}:=\mathbf{D}-\mathbf{C A}^{-1} \mathbf{B}
$$

The latter is used for Cholesky factorization $\left(\mathbf{M} \in \mathbb{H}_{n}\right.$ and $\left.\mathbf{A} \in \mathbb{F}^{1 \times 1}\right)$.

What if $\mathbf{D}$ of $\mathbf{A}$ are singular or not square?

## Generalized Schur complement

Let $\mathbf{M} \in \mathbb{F}^{m \times n}$ and

$$
\boldsymbol{\alpha}=\left(\alpha_{1}, \ldots, \alpha_{k}\right) \subseteq \llbracket m \rrbracket \quad \text { and } \quad \boldsymbol{\alpha}^{c}=\llbracket m \rrbracket \backslash \boldsymbol{\alpha}
$$

and

$$
\boldsymbol{\beta}=\left(\beta_{1}, \ldots, \beta_{\ell}\right) \subseteq \llbracket n \rrbracket \quad \text { and } \quad \boldsymbol{\beta}^{c}=\llbracket n \rrbracket \backslash \boldsymbol{\beta} .
$$

We denote

$$
\mathbf{M}[\gamma, \boldsymbol{\delta}]
$$

the $(\boldsymbol{\gamma}, \boldsymbol{\delta})$-block in $\mathbf{M}$.
The Schur complement of $\mathbf{M}[\boldsymbol{\alpha}, \boldsymbol{\beta}]$ in $\mathbf{M}$ is

$$
\mathbf{M} / \mathbf{M}[\boldsymbol{\alpha}, \boldsymbol{\beta}]=\mathbf{M}\left[\boldsymbol{\alpha}^{c}, \boldsymbol{\beta}^{c}\right]-\mathbf{M}\left[\boldsymbol{\alpha}^{c}, \boldsymbol{\beta}\right](\mathbf{M}[\boldsymbol{\alpha}, \boldsymbol{\beta}])^{\dagger} \mathbf{M}\left[\boldsymbol{\alpha}, \boldsymbol{\beta}^{c}\right]
$$

F. Zhang, The Schur complement and its applications, Vol. 4 of Numerical Methods and Algorithms, Springer-Verlag, New York.

## Approximation in the spectral norm

We will mostly establish spectral norm errors.

- Suppose $\mathbf{A} \in \mathbb{F}^{m \times n}$ and $\hat{\mathbf{A}} \in \mathbb{F}^{m \times n}$ is an approximation

$$
\|\mathbf{A}-\hat{\mathbf{A}}\| \leq \varepsilon
$$

then

1. $|\langle\mathbf{F}, \mathbf{A}\rangle-\langle\mathbf{F}, \hat{\mathbf{A}}\rangle| \leq \varepsilon\|\mathbf{F}\|_{*}$ for every matrix $\mathbf{F} \in \mathbb{F}^{m \times n}$
2. $\left|\sigma_{j}(\mathbf{A})-\sigma_{j}(\hat{\mathbf{A}})\right| \leq \varepsilon, \forall j$

How do spectral norm errors compare with Frobenius norm error measures?

