# Multi-linear Algebra <br> - Tensor Train Decomposition Lecture 18 

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29/03/2024

Recall

## Recall

- CP decomposition: Let $\mathbf{A} \in \mathbb{R}^{n_{1} \times \ldots \times n_{d}}$. Then

$$
\mathbf{A}=\sum_{p=1}^{r} \bigotimes_{i=1}^{d} \mathbf{v}_{i, p}
$$

Storage of CP format:
CP rank:

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$$

Storage of CP format: $\mathcal{O}(r n d)$
CP rank: minimal $r$ s.t. we can express $\mathbf{A}$ in the above format

## Recall

- Tucker decomposition: Let $\mathbf{A} \in \mathbb{R}^{n_{1} \times \ldots \times n_{d}}$. Then

$$
\begin{aligned}
\mathbf{A} & =\sum_{i_{1}=1}^{r_{1}} \cdots \sum_{i_{d}=1}^{r_{d}} \mathbf{C}\left[i_{1}, \ldots, i_{d}\right] \cdot \mathbf{u}_{1, i_{1}} \otimes \mathbf{u}_{2, i_{2}} \otimes \cdots \otimes \mathbf{u}_{d, i_{1}} \\
& =\mathbf{C} *_{1} \mathbf{U}_{1} *_{2} \mathbf{U}_{2} \ldots *_{d} \mathbf{U}_{d}
\end{aligned}
$$

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- Tucker decomposition: Let $\mathbf{A} \in \mathbb{R}^{n_{1} \times \ldots \times n_{d}}$. Then


Storage of Tucker format:
T-rank:

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Storage of Tucker format: $\mathcal{O}\left(r^{d}+r n d\right)$
T-rank: $\mathbf{r}=\left(r_{1}, \ldots, r_{d}\right)$
Advantage:

## Recall

- Tucker decomposition: Let $\mathbf{A} \in \mathbb{R}^{n_{1} \times \ldots \times n_{d}}$. Then


Storage of Tucker format: $\mathcal{O}\left(r^{d}+r n d\right)$
Advantage:
1 Can be computed using HOSVD
2 Closed set of low-rank tensors
3 Manifold structure on the set of tensors with fixed rank
4 Can be sketched

## Sketching Tucker

Compute a low-Tucker rank approximation using:

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The computational complexity of the most common operations scales linearly in the order if all operands are given in TT representation


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Big benefit:
The computational complexity of the most common operations scales linearly in the order if all operands are given in TT representation
$\Rightarrow$ TT decomposition unifies advantages of CP and Tucker


## TT decomposition - A tail of subspaces!

Let $\mathbf{A} \in \mathbb{R}^{n_{1} \times \ldots \times n_{d}}$ be our target tensor

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- For TT: find a hierarchy of nested subspaces

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U_{1} \subseteq \mathbb{R}^{n_{1}}, U_{2} \subseteq \mathbb{R}^{n_{1} \times n_{2}}, \ldots, U_{d-1} \subseteq \mathbb{R}^{n_{1} \times \ldots \times n_{d-1}}
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s.t. the final subspace contains the target tensor:

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with $\mathbf{A} \in U_{1} \otimes \mathbb{R}^{n_{2} \times \ldots \times n_{d}}$ with $\mathbf{A} \in U_{2} \otimes \mathbb{R}^{n_{3} \times \ldots \times n_{d}}$

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$$

$$
U_{d-1} \subseteq U_{d-2} \otimes \mathbb{R}^{n_{d-1}} \subseteq \mathbb{R}^{n_{1} \times \ldots \times n_{d-1}}
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with $\mathbf{A} \in U_{1} \otimes \mathbb{R}^{n_{2} \times \ldots \times n_{d}}$ with $\mathbf{A} \in U_{2} \otimes \mathbb{R}^{n_{3} \times \ldots \times n_{d}}$ with $\mathbf{A} \in U_{3} \otimes \mathbb{R}^{n_{4} \times \ldots \times n_{d}}$
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- Note that $U_{k} \subseteq U_{k-1} \otimes \mathbb{R}^{n_{k}} \subseteq \mathbb{R}^{n_{1} \times \ldots \times n_{k}}$ ensures that

$$
\mathbf{V}_{k, j}=\sum_{i=1}^{r_{k-1}} \mathbf{V}_{k-1, i} \otimes \mathbf{u}_{k, i, j}
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for some $\mathbf{u}_{k, i, j} \in \mathbb{R}^{n_{k}}$.

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- Writing the (orthogonal) basis as a tensor

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\mathbf{W}_{k}\left[i_{1}, \ldots, i_{k}, j\right]=\mathbf{V}_{k, j}\left[i_{1}, \ldots, i_{k}\right]
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and defining

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\mathbf{U}_{k}[i, \ell, j]=\mathbf{u}_{k, i, j}[\ell]
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with $\mathbf{W} \in \mathbb{R}^{n_{1} \times \ldots \times n_{k} \times r_{k}}$ and $\mathbf{U}_{k} \in \mathbb{R}^{r_{k-1} \times n_{k} \times r_{k}}$

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& =\sum_{i=1}^{r_{k-1}} \mathbf{W}_{k-1}\left[i_{1}, \ldots, i_{k-1}, i\right] \mathbf{U}_{k}\left[i, i_{k}, j\right] \\
& =\left(\mathbf{W}_{k-1} *_{(k),(1)} \mathbf{U}_{k}\right)\left[i_{1}, \ldots, i_{k-1}, i_{k}, j\right]
\end{aligned}
$$

## TT decomposition - A tail of subspaces!

So, recursively applied, this yields

$$
\mathbf{A}=\mathbf{W}_{d-1} *_{(d),(1)} \mathbf{U}_{d}
$$

## TT decomposition - A tail of subspaces!

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\begin{aligned}
\mathbf{A} & =\mathbf{W}_{d-1} *(d),(1) \\
& \mathbf{U}_{d} \\
& =\left(\mathbf{W}_{d-2} *_{(d-1),(1)} \mathbf{U}_{d-1}\right) *_{(d),(1)} \mathbf{U}_{d}
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& =\left(\mathbf{W}_{d-3} *_{(d-2),(1)} \mathbf{U}_{d-2}\right) *_{(d-1),(1)} \mathbf{U}_{d-1} *_{(3),(1)} \mathbf{U}_{d}
\end{aligned}
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& =\mathbf{U}_{1} *_{(3),(1)} \mathbf{U}_{2} *_{(3),(1)} \cdots *_{(3),(1)} \mathbf{U}_{d}
\end{aligned}
$$

Or as a diagram


## TT-SVD (another variant of HOSVD)

$$
\mathbf{A}_{n_{1}, \ldots, n_{d}}
$$

## TT-SVD (another variant of HOSVD)

$$
\begin{aligned}
& \mathbf{A}_{n_{1}, \ldots, n_{d}} \\
& =\mathbf{A}_{n_{2} \cdots n_{d}}
\end{aligned}
$$

$$
\text { reshape to } n_{1} \times \prod_{j \neq i} n_{j}
$$

## TT-SVD (another variant of HOSVD)

$$
\begin{aligned}
& \mathbf{A}_{n_{1}, \ldots, n_{d}} \\
& =\mathbf{A}_{n_{2} \cdots n_{d}}^{n_{1}} \\
& =\left(\mathbf{U}_{1}\right)_{r_{1}}^{n_{1}}\left(\boldsymbol{\Sigma}_{1} \mathbf{V}_{1}^{T}\right)_{n_{2} \ldots n_{d}}^{r_{1}}
\end{aligned}
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SVD

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& =\left(\mathbf{U}_{1}\right)_{r_{1}}^{n_{1}}\left(\boldsymbol{\Sigma}_{1} \mathbf{V}_{1}^{\top}\right)_{n_{3} \cdots n_{d}}^{r_{1} \cdot n_{2}}
\end{aligned}
$$

## reshape to $n_{1} \times \prod_{j \neq i} n_{j}$

SVD
reshape of $\left(\boldsymbol{\Sigma}_{1} \mathbf{V}_{1}^{\top}\right)$

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\end{aligned}
$$

SVD
reshape of $\left(\boldsymbol{\Sigma}_{1} \mathbf{V}_{1}^{\top}\right)$
$\operatorname{SVD}$ of $\left(\boldsymbol{\Sigma}_{1} \mathbf{V}_{1}^{\top}\right)$
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& \text { reshape to } n_{1} \times \prod_{j \neq i} n_{j} \\
& \text { SVD } \\
& \text { reshape of }\left(\boldsymbol{\Sigma}_{1} \mathbf{V}_{1}^{\top}\right) \\
& \text { SVD of }\left(\boldsymbol{\Sigma}_{1} \mathbf{V}_{1}^{\top}\right) \\
& \text { reshape of }\left(\boldsymbol{\Sigma}_{2} \mathbf{V}_{2}^{\top}\right) \\
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& =\left(\mathbf{U}_{1}\right)_{r_{1}}^{n_{1}}\left(\boldsymbol{\Sigma}_{1} \mathbf{V}_{1}^{\top}\right)_{n_{3} \cdots n_{d}}^{r_{1} \cdot \mathbf{U}_{1}} \\
& =\left(\mathbf{U}_{1}\right)_{r_{1}}^{n_{1}}\left(\mathbf{U}_{2}\right)_{r_{2}}^{r_{1} \cdot n_{2}}\left(\boldsymbol{\Sigma}_{2} \mathbf{V}_{2}^{\top}\right)_{n_{3} \cdots n_{d}}^{r_{2}} \\
& =\left(\mathbf{U}_{1}\right)_{r_{1}}^{n_{1}}\left(\mathbf{U}_{2}\right)_{r_{2}}^{r_{1} \cdot n_{2}}\left(\boldsymbol{\Sigma}_{2} \mathbf{V}_{2}^{\top}\right)_{n_{4} \cdots n_{d}}^{r_{2} \cdot n_{3}} \\
& =\left(\mathbf{U}_{1}\right)_{r_{1}}^{n_{1}}\left(\mathbf{U}_{2}\right)_{r_{2}}^{r_{1} \cdot n_{2}}\left(\mathbf{U}_{3}\right)_{r_{3}}^{r_{2} \cdot n_{3}}\left(\mathbf{\Sigma}_{3} \mathbf{V}_{3}^{\top}\right)_{n_{4} \cdots n_{d}}^{r_{3}} \\
& =\underbrace{\left(\mathbf{U}_{1}\right)_{r_{1}}^{n_{1}}}_{\mathbf{U}_{1}\left[n_{d}\right]} \cdots \underbrace{\left(\mathbf{U}_{d-1}\right)_{r_{d-1}}^{r_{d-2} \cdot n_{d-1}}}_{=: \mathbf{U}_{d-1}\left[n_{d-1}\right]} \underbrace{\left(\mathbf{\Sigma}_{d-1} \mathbf{V}_{d-1}^{\top}\right)_{n_{d}}^{r_{d-1}}}_{=: \mathbf{U}_{d}\left[n_{d}\right]}
\end{aligned}
$$

## TT-SVD (diagrammatically)

$$
-\overbrace{n_{1}}^{\mathcal{X}_{0} \overbrace{n_{4}} n_{n_{3}}}
$$

## TT-SVD (diagrammatically)


$n_{1} \boldsymbol{X}_{0}^{(1)}{ }_{n} n_{3} n_{4}$

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## TT-SVD (diagrammatically)



## TT-SVD (diagrammatically)



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## TT-SVD (diagrammatically)



## TT-SVD (diagrammatically)



## TT-SVD (diagrammatically)



## TT-SVD (diagrammatically)



## TT-SVD (diagrammatically)



## TT-SVD (diagrammatically)



## TT Pseudo-code

We introduce the notation

$$
\mathbf{A}^{\langle k\rangle}=\operatorname{MAT}_{(1, \ldots, k)}(\mathbf{A})
$$

for the matricization that flattens the first $k$ and the last $d-k$ modes
Algorithm:
Input: Target tensor A, target $\operatorname{rank}\left(r_{1}, \ldots, r_{d}\right)$ Output: Component tensors $\mathbf{U}_{i} \in \mathbb{R}^{r_{i-1} \times n_{i} \times r_{i}}$

$$
\begin{aligned}
& \mathbf{A}_{1}=\mathbf{A}^{\langle 1\rangle} \\
& \tilde{\mathbf{U}}_{1}, \boldsymbol{\Sigma}_{1}, \mathbf{V}_{1}=\operatorname{SVD}\left(\mathbf{A}_{1}, r_{1}\right) \\
& \mathbf{U}_{1}=\operatorname{unfild}\left(\tilde{\mathbf{U}}_{1}\right) \\
& \mathcal{V}_{1}=\operatorname{unfold}\left(\boldsymbol{\Sigma}_{1} \mathbf{V}_{1}^{\top}\right) \\
& A_{2}=\mathcal{V}_{1}^{\langle 2\rangle} \\
& \text { for } k=2: d-1 \\
& \quad \tilde{\mathbf{U}}_{k}, \boldsymbol{\Sigma}_{k}, \mathbf{V}_{k}=\operatorname{SVD}\left(\mathbf{A}_{k}, r_{k}\right) \\
& \quad \mathbf{U}_{k}=\operatorname{unfold}\left(\tilde{\mathbf{U}}_{k}\right) \\
& \quad \mathcal{V}_{k}=\operatorname{unfold}\left(\boldsymbol{\Sigma}_{k} \mathbf{V}_{k}^{\top}\right) \\
& \quad A_{k+1}=\mathcal{V}_{k}^{\langle 2\rangle} \\
& \tilde{\mathbf{U}}_{d}=\mathcal{V}_{d-1} \\
& \mathbf{U}_{d}=\operatorname{unfold}^{\left(\tilde{\mathbf{U}}_{d}\right)}
\end{aligned}
$$

## TT decomposition

Let $\mathbf{A} \in \mathbb{R}^{n_{1} \times \ldots \times n_{d}}$. We call a factorization

$$
\mathbf{A}\left[i_{1}, \ldots, i_{d}\right]=\mathbf{U}_{1}\left[i_{1}\right] \mathbf{U}_{2}\left[i_{2}\right] \cdots \mathbf{U}_{d}\left[i_{d}\right]
$$

a TT representation of $\mathbf{A}$.
Note: This reveals the alternative name Matrix-product-states

Storing in TT format scales as

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a TT representation of $\mathbf{A}$.
Note: This reveals the alternative name Matrix-product-states

Storing in TT format scales as

$$
\mathcal{O}\left(r^{2} d n\right)
$$

## Accessing Entries

Let consider $\mathbf{A} \in \mathbb{R}^{n_{1} \times n_{2} \times n_{3} \times n_{4}}$. Then

$$
\mathbf{A}\left[i_{1}, \ldots, i_{4}\right]=\sum_{k_{3}=1}^{r_{3}} \sum_{k_{2}=1}^{r_{2}} \sum_{k_{1}=1}^{r_{1}} \mathbf{U}_{1}\left[1, i_{1}, k_{1}\right] \mathbf{U}_{2}\left[k_{1}, i_{2}, k_{2}\right] \mathbf{U}_{3}\left[k_{2}, i_{3}, k_{3}\right] \mathbf{U}_{4}\left[k_{3}, i_{4}, 1\right]
$$

## Accessing Entries

Let consider $\mathbf{A} \in \mathbb{R}^{n_{1} \times n_{2} \times n_{3} \times n_{4}}$. Then

$$
\begin{aligned}
& \mathbf{A}\left[i_{1}, \ldots, i_{4}\right]=\sum_{k_{3}=1}^{r_{3}} \sum_{k_{2}=1}^{r_{2}} \sum_{k_{1}=1}^{r_{1}} \mathbf{U}_{1}\left[1, i_{1}, k_{1}\right] \mathbf{U}_{2}\left[k_{1}, i_{2}, k_{2}\right] \mathbf{U}_{3}\left[k_{2}, i_{3}, k_{3}\right] \mathbf{U}_{4}\left[k_{3}, i_{4}, 1\right] \\
&=\sum_{k_{3}=1}^{r_{3}} \sum_{k_{2}=1}^{r_{2}} \underbrace{\left(\sum_{k_{1}=1}^{r_{1}} \mathbf{U}_{1}\left[1, i_{1}, k_{1}\right] \mathbf{U}_{2}\left[k_{1}, i_{2}, k_{2}\right]\right.}_{\substack{\in \mathcal{O}(r) \text { for fixed } i_{1}, i_{2}, k_{2} \\
\in \mathcal{O}\left(r^{2}\right) \text { since } k_{2} \in \llbracket r_{2} \rrbracket}}) \\
& \mathbf{U}_{3}\left[k_{2}, i_{3}, k_{3}\right] \mathbf{U}_{4}\left[k_{3}, i_{4}, 1\right]
\end{aligned}
$$

## Accessing Entries

Let consider $\mathbf{A} \in \mathbb{R}^{n_{1} \times n_{2} \times n_{3} \times n_{4}}$. Then

$$
\begin{aligned}
\mathbf{A}\left[i_{1}, \ldots, i_{4}\right] & =\sum_{k_{3}=1}^{r_{3}} \sum_{k_{2}=1}^{r_{2}} \sum_{k_{1}=1}^{r_{1}} \mathbf{U}_{1}\left[1, i_{1}, k_{1}\right] \mathbf{U}_{2}\left[k_{1}, i_{2}, k_{2}\right] \mathbf{U}_{3}\left[k_{2}, i_{3}, k_{3}\right] \mathbf{U}_{4}\left[k_{3}, i_{4}, 1\right] \\
& =\sum_{k_{3}=1}^{r_{3}} \sum_{k_{2}=1}^{\sum_{2}} \underbrace{\left(\sum_{k_{1}=1}^{r_{1}} \mathbf{U}_{1}\left[1, i_{1}, k_{1}\right] \mathbf{U}_{2}\left[k_{1}, i_{2}, k_{2}\right]\right)}_{\substack{\left.\in \mathcal{O}(r) \text { for fixed } i_{1}, i_{2}, k_{2} \\
\in \mathcal{O}\left(r^{2}\right) \text { since } k_{2} \in \llbracket r_{2}\right]}} \mathbf{U}_{3}\left[k_{2}, i_{3}, k_{3}\right] \mathbf{U}_{4}\left[k_{3}, i_{4}, 1\right] \\
& =\sum_{k_{3}=1}^{r_{3}} \underbrace{\left(\sum_{k_{2}=1}^{r_{2}} \mathbf{W}_{1}\left[1, i_{1}, i_{2}, k_{2}\right] \mathbf{U}_{3}\left[k_{2}, i_{3}, k_{3}\right]\right)}_{\in \mathcal{O}\left(r^{2}\right)} \mathbf{U}_{4}\left[k_{3}, i_{4}, 1\right] \\
& =\ldots
\end{aligned}
$$

Generalizing this idea leads to the computational scaling $\mathcal{O}\left(d r^{2}\right)$

## Adding TT decomposition

$$
\begin{aligned}
& (\mathbf{A}+\overline{\mathbf{A}})\left[i_{1}, \ldots, i_{d}\right] \\
& =\mathbf{U}_{1, i_{1}} \mathbf{U}_{2, i_{2}} \ldots \mathbf{U}_{d-1, i_{d-1}} \mathbf{U}_{d, i_{d}}+\overline{\mathbf{U}}_{1, i_{1}} \overline{\mathbf{U}}_{2, i_{2}} \ldots \overline{\mathbf{U}}_{d-1, i_{d-1}} \overline{\mathbf{U}}_{d, i_{d}} \\
& =\left(\begin{array}{ll}
\mathbf{U}_{1, i_{1}} & \overline{\mathbf{U}}_{1, i_{1}}
\end{array}\right)\left(\begin{array}{cc}
\mathbf{U}_{2, i_{2}} & \mathbf{0} \\
\mathbf{0} & \overline{\mathbf{U}}_{2, i_{2}}
\end{array}\right) \cdots\left(\begin{array}{cc}
\mathbf{U}_{d-1, i_{d-1}} & \mathbf{0} \\
\mathbf{0} & \overline{\mathbf{U}}_{d-1, i_{d-1}}
\end{array}\right)\binom{\mathbf{U}_{d, i_{d}}}{\overline{\mathbf{U}}_{d, i_{d}}} \\
& =\mathbf{W}_{1, i_{1}} \mathbf{W}_{2, i_{2}} \ldots \mathbf{W}_{d-1, i_{d-1}} \mathbf{W}_{d, i_{d}}
\end{aligned}
$$

Which is a valid TT representation of $\mathbf{A}+\overline{\mathbf{A}}$ of $\operatorname{rant} \mathbf{r}+\overline{\mathbf{r}}$. Thus the scaling is

$$
\mathcal{O}\left(d(r+\bar{r})^{2}\right)
$$

where $r=\max (\mathbf{r})$, and $\bar{r}=\max (\overline{\mathbf{r}})$

## Other operations

| Operation | TT | Tucker | CP |
| :---: | :---: | :---: | :---: |
| Had. Prod. | $\mathcal{O}\left(n d r^{2} \bar{r}^{2}\right)$ | $\mathcal{O}\left(n d r \bar{r}+r^{d} \bar{r}^{d}\right)$ | $\mathcal{O}(n d r \bar{r})$ |
| Frob. In. Prod. | $\mathcal{O}\left(n d r^{3}\right)$ | $\mathcal{O}\left(n d r \bar{r}+d r \bar{r}^{d}+r^{d}\right)$ | $\mathcal{O}(n d r \bar{r})$ |
| Frob. Norm | $\mathcal{O}\left(r^{2} n\right)$ | $\mathcal{O}\left(r^{d}\right)$ | $\mathcal{O}\left(n d r^{2}\right)$ |
| $k$-mode Prod. | $\mathcal{O}\left(m n r^{2}\right)$ | $\mathcal{O}\left(m n r+m r^{2}+r^{d+1}\right)$ | $\mathcal{O}((d+m) n r)$ |

