



Fast & Forward Stable Randomized Algorithm for LLSP

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Linear Least Square Problem



The overdetermined linear least squares problem is given

$$\mathbf{x} = \arg \min_{\mathbf{y} \in \mathbb{R}^n} \|\mathbf{b} - \mathbf{A}\mathbf{y}\|$$

for $\mathbf{A} \in \mathbb{R}^{m \times n}$, $\mathbf{b} \in \mathbb{R}^m$





Algorithm - Iterative Sketching



Input: Matrix $\mathbf{A} \in \mathbb{R}^{m \times n}$, RHS $\mathbf{b} \in \mathbb{R}^m$, Embedding dimension d iteration count q .

Output: Approx solution x_q to least squares.

(1) $\mathbf{S} \leftarrow d \times m$

(1) Sketching matrix

(2) $\mathbf{B} \leftarrow \mathbf{S}\mathbf{A}$

(2) Applying the sketch

(3) $(\mathbf{Q}, \mathbf{R}) \leftarrow \text{QR}(\mathbf{B}, \text{'econ'})$

(4) $\mathbf{x}_0 \leftarrow \mathbf{R}^{-1}(\mathbf{Q}^*(\mathbf{S}\mathbf{b}))$

(3,4) Householder QR

(5) for $i = 0 : q - 1$

(6) $\mathbf{x}_{i+1} \leftarrow \mathbf{x}_i + \mathbf{R}^{-1}\mathbf{R}^{-\top}\mathbf{A}^{\top}(\mathbf{b} - \mathbf{A}\mathbf{x}_i)$ (6) Update rule

(7) end





Algorithm - Justification



$$(1) \quad (\mathbf{A}^T \mathbf{A})\mathbf{x} = \mathbf{A}^T \mathbf{b}$$

$$(2) \quad (\mathbf{S}\mathbf{A})^T(\mathbf{S}\mathbf{A}) \approx \mathbf{A}^T \mathbf{A}$$

$$(3) \quad \mathbf{A}^T \mathbf{A}(\mathbf{x} - \mathbf{x}_i) = \mathbf{A}^T(\mathbf{b} - \mathbf{A}\mathbf{x}_i)$$

$$(4) \quad (\mathbf{S}\mathbf{A})^T(\mathbf{S}\mathbf{A})\mathbf{d}_i = \mathbf{A}^T(\mathbf{b} - \mathbf{A}\mathbf{x}_i)$$

$$(5) \quad \text{where, } \mathbf{x} \approx \mathbf{x}_{i+1} = \mathbf{x}_i + \mathbf{d}_i$$

$$(6) \quad \mathbf{S}\mathbf{A} = \mathbf{Q}\mathbf{R}$$

$$(7) \quad (\mathbf{S}\mathbf{A})^T(\mathbf{S}\mathbf{A}) = \mathbf{R}^T \mathbf{R}$$

$$(8) \quad \mathbf{d}_i = \mathbf{R}^{-1} \mathbf{R}^{-T} \mathbf{A}^T(\mathbf{b} - \mathbf{A}\mathbf{x}_i)$$

$$(9) \quad \mathbf{x}_{i+1} = \mathbf{x}_i + \mathbf{R}^{-1} \mathbf{R}^{-T} \mathbf{A}^T(\mathbf{b} - \mathbf{A}\mathbf{x}_i)$$





Notational Notes



♥ **Relative Forward Error:** $FE(\hat{\mathbf{x}}) = \frac{\|\mathbf{x} - \hat{\mathbf{x}}\|}{\|\mathbf{x}\|}$

♥ **Relative Residual:** $RE(\hat{\mathbf{x}}) = \frac{\|\mathbf{r}(\mathbf{x}) - \mathbf{r}(\hat{\mathbf{x}})\|}{\|\mathbf{r}(\mathbf{x})\|}$ where $\mathbf{r}(\mathbf{y}) = \mathbf{b} - \mathbf{A}\mathbf{y}$

♥ **Relative Backward Error:** $BE(\hat{\mathbf{x}}) = \min \left\{ \frac{\|\Delta \mathbf{A}\|_F}{\|\mathbf{A}\|_F} : \hat{\mathbf{x}} = \arg \min_{\mathbf{v} \in \mathbb{R}^n} \|\mathbf{b} - (\mathbf{A} + \Delta \mathbf{A})\mathbf{v}\| \right\}$

♥ **Machine Epsilon:** $u = \epsilon_{mach} \approx 10^{-16}$

♥ **Condition Number:** $\kappa(\mathbf{A}) = \kappa = \|\mathbf{A}\| \|\mathbf{A}^{-1}\| = \sigma_{\max}(\mathbf{A}) / \sigma_{\min}(\mathbf{A})$

If BE is small, then computed solution is the *true solution to the nearly the right least-square problem*
Informally, if $BE \approx u$ then the algorithm is *backwards stable*





Comparing with Backward Stable Algo



- ♥ Householder QR is *backwards stable* i.e. $\hat{\mathbf{x}} = \arg \min_{\mathbf{y} \in \mathbb{R}^n} \|(\mathbf{b} + \Delta \mathbf{b}) - (\mathbf{A} + \Delta \mathbf{A})\mathbf{y}\|$ with perturbations of size $\|\Delta \mathbf{A}\| \leq cu\|\mathbf{A}\|$ and $\|\Delta \mathbf{b}\| \leq cu\|\mathbf{b}\|$ provided $cu < 1$
- ♥ **Main Result:** iterative sketching has forward errors and residual errors comparable to a backward stable algorithm (e.g. Householder QR)

But what are the forward and residual errors for a backward stable algorithm?



Wedin's Perturbation Theorem

Let $\mathbf{A}, \Delta\mathbf{A} \in \mathbb{R}^{m \times n}$ and $\mathbf{b}, \Delta\mathbf{b} \in \mathbb{R}^m$ and suppose that $\|\Delta\mathbf{A}\| \leq \epsilon\|\mathbf{A}\|$ and $\|\Delta\mathbf{b}\| \leq \epsilon\|\mathbf{b}\|$ for some $(0, 1)$. Then set

$$\mathbf{x} = \arg \min_{\mathbf{y} \in \mathbb{R}^n} \|\mathbf{b} - \mathbf{A}\mathbf{y}\|$$

$$\hat{\mathbf{x}} = \arg \min_{\mathbf{y} \in \mathbb{R}^n} \|(\mathbf{b} + \Delta\mathbf{b}) - (\mathbf{A} + \Delta\mathbf{A})\mathbf{y}\|$$

Then if $\epsilon\kappa(\mathbf{A}) \leq 0.1$

$$\|\mathbf{x} - \hat{\mathbf{x}}\| \leq 2.23\kappa(\mathbf{A}) \left(\|\mathbf{x}\| + \frac{\kappa(\mathbf{A})}{\|\mathbf{A}\|} \|\mathbf{r}(\mathbf{x})\| \right) \epsilon$$
$$\|\mathbf{r}(\mathbf{x}) - \mathbf{r}(\hat{\mathbf{x}})\| \leq 2.23 (\|\mathbf{A}\| \|\mathbf{x}\| + \kappa(\mathbf{A}) \|\mathbf{r}(\mathbf{x})\|) \epsilon$$



Results from Wedin's Theorem



Assuming κu sufficiently small, then the solution $\hat{\mathbf{x}}$ to a *backwards stable* least square solver satisfies

$$\begin{aligned}\|\mathbf{x} - \hat{\mathbf{x}}\| &\leq c\kappa(\mathbf{A}) \left(\|\mathbf{x}\| + \frac{\kappa}{\|\mathbf{A}\|} \|\mathbf{r}(\mathbf{x})\| \right) u \\ \|\mathbf{r}(\mathbf{x}) - \mathbf{r}\hat{\mathbf{x}}\| &\leq c' (\|\mathbf{A}\| \|\mathbf{x}\| + \kappa \|\mathbf{r}(\mathbf{x})\|) u\end{aligned}$$

An algorithm is said to be *forward stable* if the computed solution satisfies bounds of this form



Theorem VI: Iterative sketching is forward stable

Let $\mathbf{S} \in \mathbb{R}^{d \times m}$ be a subspace embedding for $\text{range}([\mathbf{A} \ \mathbf{b}])$ with distortion $\epsilon \in (0, 0.29]$. There exists a constant $c_1 > 0$ which depends polynomially on m, n , and d such that $c_1 \kappa u < 1$ and multiplication by \mathbf{S} is forward stable, then the numerically computed iterates $\hat{\mathbf{x}}_i$ by Iterative Method satisfy bounds

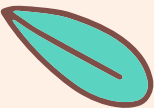
$$\begin{aligned}\|\mathbf{x} - \hat{\mathbf{x}}_i\| &\leq 20\sqrt{\epsilon}\kappa(g_{IS} + c_1\kappa u)^i \frac{\|\mathbf{r}(\mathbf{x})\|}{\|\mathbf{A}\|} + c_1\kappa u \left[\|\mathbf{x}\| + \frac{\kappa}{\|\mathbf{A}\|} \|\mathbf{r}(\mathbf{x})\| \right] \\ \|\mathbf{r}(\mathbf{x}) - \mathbf{r}(\hat{\mathbf{x}}_i)\| &\leq 20\sqrt{\epsilon}(g_{IS} + c_1\kappa u)^i \|\mathbf{r}(\mathbf{x})\| + c_1 u [\|\mathbf{A}\| \|\mathbf{x}\| + \kappa \|\mathbf{r}(\mathbf{x})\|]\end{aligned}$$

In particular, if $g_{IS} + c_i \kappa u \leq 0.9$ then the iterative sketching produces a solution $\hat{\mathbf{x}}$ such it satisfies the forward stability error bounds





Informally...

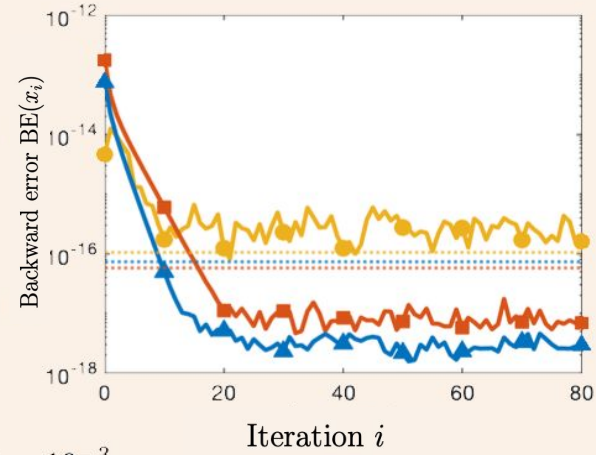
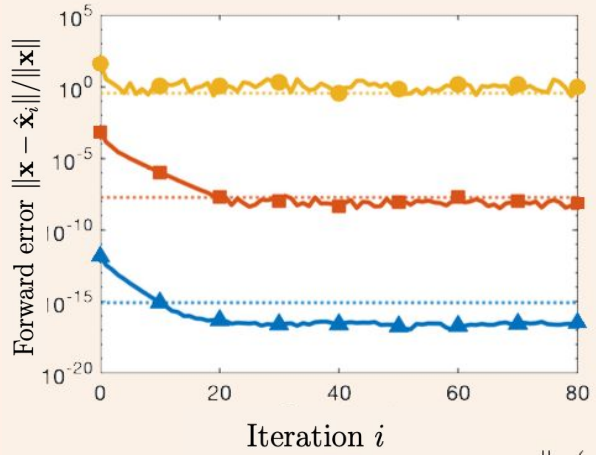


(Informal Theorem I) In FPA, iterative sketching produces iteratives $\hat{\mathbf{x}}_0, \hat{\mathbf{x}}_1, \dots$ with errors $\|\mathbf{x} - \hat{\mathbf{x}}_i\|$ and residual errors $\|\mathbf{r}(\mathbf{x}) - \mathbf{r}(\hat{\mathbf{x}}_i)\|$ that converge geometrically until they reach roughly the same accuracy as Householder QR.

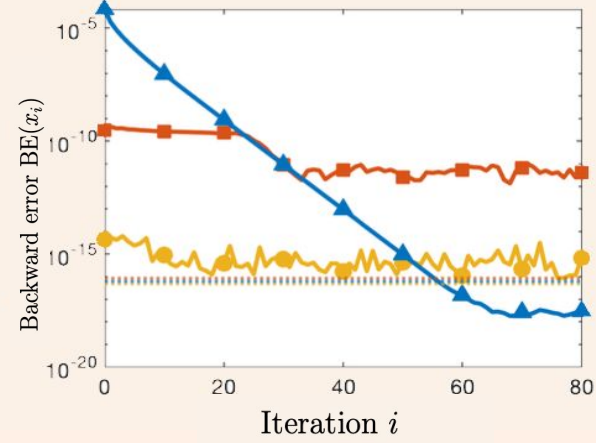
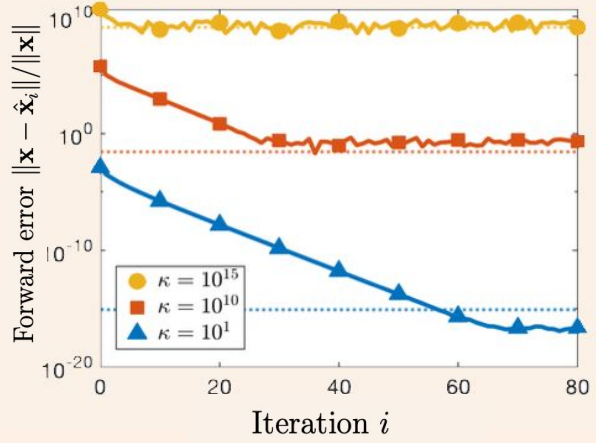




$$\|r(x)\| = 10^{-12}$$



$$\|r(x)\| = 10^{-3}$$





Convergence - Theorem V

Let \mathbf{S} be a subspace embedding with distortion $0 < \epsilon < 1 - \frac{1}{\sqrt{2}}$. Then the \mathbf{x}_i satisfy the following bounds,

$$(1) \quad \|\mathbf{x} - \mathbf{x}_i\| < (8 - 2\sqrt{2})\sqrt{\epsilon}\kappa g_{IS}^i \frac{\|\mathbf{r}(\mathbf{x})\|}{\|\mathbf{A}\|} \quad (2) \quad \|\mathbf{r}(\mathbf{x}) - \mathbf{r}(\mathbf{x}_i)\| < (8 - 2\sqrt{2})\sqrt{\epsilon}g_{IS}^i \|\mathbf{r}(\mathbf{x})\|$$

The convergence rate g_{IS} is,

$$g_{IS} = \frac{\epsilon(2 - \epsilon)}{(1 - \epsilon)^2} \leq (2 + \sqrt{2})\epsilon$$



Lemma: Singular Value Bounds

Let \mathbf{S} be a subspace embedding for $\text{range}(\mathbf{A})$ with distortion $\epsilon \in (0, 1)$ and $\mathbf{QR} = \mathbf{SA}$ be a reduced QR decomposition of \mathbf{SA} . Then \mathbf{R} satisfies the bounds

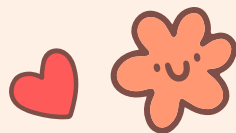
$$\sigma_{\max}(\mathbf{R}) \leq (1 + \epsilon)\sigma_{\max}(\mathbf{A})$$

$$\sigma_{\min}(\mathbf{R}) \geq (1 - \epsilon)\sigma_{\min}(\mathbf{A})$$

In addition, \mathbf{AR}^{-1} satisfies the bounds

$$\sigma_{\max}(\mathbf{AR}^{-1}) \leq \frac{1}{1 - \epsilon}$$

$$\sigma_{\min}(\mathbf{AR}^{-1}) \geq \frac{1}{1 + \epsilon}$$



Proof of Lemma

$$\sigma_{\max}(\mathbf{R}) \leq (1 + \epsilon)\sigma_{\max}(\mathbf{A})$$

Recall, $\mathbf{SA} = \mathbf{QR}$ and \mathbf{S} is a subspace embedding for subspace V such that $(1 - \epsilon)\|\mathbf{v}\| \leq \|\mathbf{Sv}\| \leq (1 + \epsilon)\|\mathbf{v}\|$ for $\epsilon \in (0, 1)$ and $\mathbf{v} \in V$.

As \mathbf{Q} unitary, \mathbf{QR} and \mathbf{R} have the same singular values. Establishing the upper bound,

$$\sigma_{\max}(\mathbf{R}) = \sigma_{\max}(\mathbf{QR}) = \sigma_{\max}(\mathbf{SA}) \quad (1)$$

$$\sigma_{\max}(\mathbf{SA}) = \|\mathbf{SA}\| \quad (2)$$

$$= \max_{\|\mathbf{v}\|=1} \|\mathbf{SAv}\| \quad (3)$$

$$\leq (1 + \epsilon) \max_{\|\mathbf{v}\|=1} \|\mathbf{Av}\| \quad (4)$$

$$= (1 + \epsilon)\sigma_{\max}(\mathbf{A}) \quad (5)$$

$$(6)$$





Proof cont.

$$\sigma_{\min}(\mathbf{R}) \geq (1 - \epsilon)\sigma_{\min}(\mathbf{A})$$



Establishing the lower bound,

$$\sigma_{\min}(\mathbf{R}) = \sigma_{\min}(\mathbf{SA}) \tag{1}$$

$$= \min_{\|\mathbf{v}\|=1} \|\mathbf{SA}\mathbf{v}\| \tag{2}$$

$$\geq (1 - \epsilon) \min_{\|\mathbf{v}\|=1} \|\mathbf{A}\mathbf{v}\| \tag{3}$$

$$= (1 - \epsilon)\sigma_{\min}(\mathbf{A}) \tag{4}$$





Proof cont.

$$\sigma_{\max}(\mathbf{A}\mathbf{R}^{-1}) \leq \frac{1}{1 - \epsilon}$$

For the bounds of $\mathbf{A}\mathbf{R}^{-1}$, establishing the upper bounds

$$\begin{aligned}\sigma_{\max}(\mathbf{A}\mathbf{R}^{-1}) &= \max_{\|\mathbf{v}\|=1} \|\mathbf{A}\mathbf{R}^{-1}\mathbf{v}\| \\ &\leq \frac{1}{1 - \epsilon} \max_{\|\mathbf{v}\|=1} \|\mathbf{S}\mathbf{A}\mathbf{R}^{-1}\mathbf{v}\| \\ &= \frac{1}{1 - \epsilon} \max_{\|\mathbf{v}\|=1} \|\mathbf{Q}\mathbf{v}\| \\ &= \frac{1}{1 - \epsilon} \sigma_{\max}(\mathbf{Q}) \\ &= \frac{1}{1 - \epsilon}\end{aligned}$$





Proof cont.

$$\sigma_{\min}(\mathbf{A}\mathbf{R}^{-1}) \geq \frac{1}{1 + \epsilon}$$

Establishing the lower bound,

$$\begin{aligned}\sigma_{\min}(\mathbf{A}\mathbf{R}^{-1}) &= \min_{\|\mathbf{v}\|=1} \|\mathbf{A}\mathbf{R}^{-1}\mathbf{v}\| \\ &\geq \frac{1}{1 + \epsilon} \min_{\|\mathbf{v}\|=1} \|\mathbf{S}\mathbf{A}\mathbf{R}^{-1}\mathbf{v}\| \\ &= \frac{1}{1 + \epsilon} \min_{\|\mathbf{v}\|=1} \|\mathbf{Q}\mathbf{v}\| \\ &= \frac{1}{1 + \epsilon} \sigma_{\min}(\mathbf{Q}) \\ &= \frac{1}{1 + \epsilon}\end{aligned}$$





Lemma: Sketch-and-Solve



In the sketch-and-solve, compute

$$\mathbf{x}_0 = \arg \min_{\mathbf{y} \in \mathbb{R}^n} \|(\mathbf{S}\mathbf{A})\mathbf{y} - \mathbf{S}\mathbf{b}\|$$

Computed numerically,

$$\mathbf{x}_0 = \mathbf{R}^{-1}(\mathbf{Q}^*(\mathbf{S}\mathbf{b}))$$

Then, we have the guarantee that

$$\|\mathbf{r}(\mathbf{x}_0)\| \leq \frac{1 + \epsilon}{1 - \epsilon} \|\mathbf{r}(\mathbf{x})\| \quad (1)$$

$$\|\mathbf{r}(\mathbf{x}) - \mathbf{r}(\mathbf{x}_0)\| \leq \frac{2\sqrt{\epsilon}}{1 - \epsilon} \|\mathbf{r}(\mathbf{x})\| \quad (2)$$

$$\|\mathbf{x} - \mathbf{x}_0\| \leq \frac{2\sqrt{\epsilon}}{1 - \epsilon} \frac{\kappa}{\|\mathbf{A}\|} \|\mathbf{r}(\mathbf{x})\| \quad (3)$$





Convergence Proof

$$(1) \mathbf{x}_{i+1} = (\mathbf{I} - \mathbf{R}^{-1}\mathbf{R}^{-T}\mathbf{A}^T\mathbf{A})\mathbf{x}_i + \mathbf{R}^{-1}\mathbf{R}^{-T}\mathbf{A}^T\mathbf{b}$$

$$(2) \mathbf{x} = (\mathbf{I} - \mathbf{R}^{-1}\mathbf{R}^{-T}\mathbf{A}^T\mathbf{A})\mathbf{x} + \mathbf{R}^{-1}\mathbf{R}^{-T}\mathbf{A}^T\mathbf{b}$$

$$(3) \mathbf{x} - \mathbf{x}_{i+1} = \mathbf{R}^{-1}\mathbf{G}\mathbf{R}(\mathbf{x} - \mathbf{x}_i) \text{ where } \mathbf{G} := \mathbf{I} - \mathbf{R}^{-T}\mathbf{A}^T\mathbf{A}\mathbf{R}^{-1}$$




$$(4) \mathbf{x} - \mathbf{x}_i = \mathbf{R}^{-1}\mathbf{G}^i\mathbf{R}(\mathbf{x} - \mathbf{x}_0)$$

$$(5) \|\mathbf{SA}(\mathbf{x} - \mathbf{x}_i)\| \leq \|\mathbf{Q}\| \|\mathbf{G}\|^i \|\mathbf{R}(\mathbf{x} - \mathbf{x}_0)\| = \|\mathbf{G}\|^i \|\mathbf{SA}(\mathbf{x} - \mathbf{x}_0)\|$$

$$(6) \|\mathbf{G}\| \leq \max \left\{ \left(\frac{1}{1-\epsilon} \right)^2 - 1, 1 - \left(\frac{1}{1+\epsilon} \right)^2 \right\} = \frac{\epsilon(2-\epsilon)}{(1-\epsilon)^2} = g_{IS}$$

$$(7) \|\mathbf{SA}(\mathbf{x} - \mathbf{x}_0)\| \leq (1 + \epsilon) \|\mathbf{A}(\mathbf{x} - \mathbf{x}_0)\| = (1 + \epsilon) \|\mathbf{r}(\mathbf{x}) - \mathbf{r}(\mathbf{x}_0)\| \leq 2 \frac{1+\epsilon}{1-\epsilon} \sqrt{\epsilon} \|\mathbf{r}(\mathbf{x})\|$$

$$(8) \|\mathbf{r}(\mathbf{x}) - \mathbf{r}(\mathbf{x}_i)\| \leq \frac{1}{1-\epsilon} \|\mathbf{SA}(\mathbf{x} - \mathbf{x}_i)\| \leq \frac{1}{1-\epsilon} \|\mathbf{G}\|^i \|\mathbf{SA}(\mathbf{x} - \mathbf{x}_0)\| \leq 2 \frac{1+\epsilon}{(1-\epsilon)^2} \sqrt{\epsilon} g_{IS}^i \|\mathbf{r}(\mathbf{x})\|$$

$$(9) \|\mathbf{r}(\mathbf{x}) - \mathbf{r}(\mathbf{x}_i)\| = \|\mathbf{A}(\mathbf{x} - \mathbf{x}_i)\| \geq \sigma_{\min}(\mathbf{A}) \|\mathbf{x} - \mathbf{x}_i\|$$




Thanks!

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